# Direct frequency-domain 3D acoustic solver with intermediate data compression benchmarked against time-domain modeling for full-waveform inversion applications 

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#### Abstract

We have developed a fast direct solver for numerical simulation of acoustic waves in 3D heterogeneous media. The Helmholtz equation is approximated by a 27 -point finite-difference stencil of second-order accuracy that is optimized to reduce the numerical dispersion. Due to the optimization, dispersion errors less than $1 \%$ are achieved with model discretization of only five points per wavelength. Wave-propagation problem requires solving a large system of linear equations with complex sparse symmetric coefficient matrix with seismic shots representing the right-hand side. The first step is the triangular factorization of the coefficient matrix followed by solving systems of linear equations with


triangular coefficient matrices as a second step. Having defined the triangular factors, the second step is very cheap, and its linear scaling with respect to the number of shots is the main advantage of direct methods. To reduce memory consumption and computational time at the factorization step, the lower triangular factors are compressed using a low-rank approximation of their nonzero blocks. The compression enables rapid solving of systems of more than $10^{8}$ equations corresponding to realistic geophysical models. Accuracy and performance comparison of our solver with a highly optimized time-domain solver proves that these approaches complement each other - depending on the problem size and computing configuration, either solver may be preferable.

## INTRODUCTION

Numerical simulation of acoustic wavefields in the frequency domain represents an important capability to solve problems arising in exploration geophysics. In particular, it serves as an engine for acoustic frequency-domain full-waveform inversion (FD FWI) (Mulder and Plessix, 2004; Shin and Cha, 2008; Virieux et al., 2009; Etienne et al., 2014). For macrovelocity reconstruction, such a simulation is usually performed many times for several low frequencies (up to 20 Hz ) at each iteration of this process.
In acoustic simulations, the pressure wavefield is excited by a point source working as a harmonic oscillator at a particular frequency. In exploration seismics, the number of shots can be in the tens of thousands or more, with receivers also numbering in the tens of thousands, especially with the proliferation of high-channel-count land
seismic acquisition systems, some of which reach one million channels (Ourabah et al., 2015).

The time- or frequency-domain approaches can be used to perform seismic modeling and FWI (Vigh and Starr, 2008; Virieux and Operto, 2009; Plessix, 2017). A conventional solver is based on a timedomain simulation followed by a Fourier transform to convert the time-domain wavefields to the frequency domain. A common strategy consists of distributing the seismic sources over processors. There is no need to exchange data between processors, and the strategy becomes very effective. Good scalability and moderate memory requirements make the approach attractive for industrial purposes. However, to be efficient in parallel processing tens of thousands of seismic sources, the approach requires significant computational resources. In this paper, we provide some experimental data on the solver developed by the Seiscope consortium.

[^0]Alternatively, seismic modeling and FWI can be performed in the frequency domain. A problem on the time harmonic solutions of the wave equation transforms into a boundary value problem for the Helmholtz equation. Absorbing boundary conditions are set to reduce the reflections of waves from the boundaries. Only a few discrete frequencies are needed to build a reliable macrovelocity model, and this is the main advantage of using the frequency-domain approach. From a practical perspective, in 3D simulations, this requires solving a system of linear equations with a huge and sparse coefficient matrix, which is some kind of approximation of the boundary value problem. The right sides of the system represent the seismic sources.

Iterative methods can be used to solve the system of linear equations. The Krylov subspace iteration methods demonstrate good efficiency (Plessix, 2007; Erlangga and Nabben, 2008; Belonosov et al., 2017, 2018) and moderate memory requirements, but they require good preconditioners and sometimes convergence may be slow or even lost.

With the increased computational power of modern computers (especially distributed memory systems, or clusters), it is now becoming more feasible to apply direct methods for solving systems of linear equations with sparse coefficient matrices. The LDLT decomposition of the coefficient matrix is first performed before computing the solutions by forward/backward substitutions. The factorization is computationally expensive, but it is performed only once per frequency because it is independent of the right-side term. The LDLT factorization leads to significant fill-in of the matrix (i.e., the number of nonzeros increases). The fill-in strongly depends on ordering the columns and rows of the matrix. To reduce the fill-in, the nested dissection (ND) method to order the matrix entries is used. For 3D problems solved with a finite-difference (FD) method, use of the ND method results in reduction of the fill-in by one order of magnitude $\left(O\left(n^{4}\right)\right.$ versus $\left.O\left(n^{5}\right)\right)$, where $n$ stands for the dimension of a 3D $n^{3}$ cubic grid (George, 1973; George et al., 1994).

A wide class of matrices arising as discretization from partial differential equations have a so-called low-rank property. That is, the fill-in in its direct factorization has low rank off-diagonal blocks. Based on this property, the fill-in can be approximated by low-rank matrices or matrices of a special hierarchical structure with lowrank blocks (Chandrasekaran et al., 2006; Xia et al., 2010; Ghysels et al., 2016). Structured methods (Xia, 2013; Wang et al., 2016) are based on using the structural low-rank approximation, and this helps to improve the theoretical performance.

In this paper, we present a multifrontal direct method of solving a system of linear equations that arises as an optimal 27-point FD approximation of a boundary value problem for the Helmholtz equation. In this method, the coefficient matrix constructed using a special kind of ordering is factorized in a product of triangular matrices. To find the unknowns, two systems with triangular coefficient matrices need to be solved.

Compared with the cost of the matrix factorization, extra right sides do not appreciably increase the computational cost (Bakulin et al., 2018). To reduce the issue of inordinate memory consumption typical for direct solvers, we apply data compression based on the low-rank approximation and hierarchical formats of storing data (Aminfar et al., 2016). We further implement message passing interface (MPI)-based parallelization inside the solver to optimize execution on distributed memory high-performance computers (clusters).

Several ways of using the low-rank property in algorithms for numerical solution of frequency-domain wave equations are known. Block low-rank methods (Amestoy et al., 2015; Pichon et al., 2018) have been proposed as simple alternatives to structured methods. Although they cannot reach the performance of the structured methods, simplicity of implementation is their attractive feature. Hierarchically semiseparable formats (Xia et al., 2010; Wang et al., 2011; Xia, 2013) have been proposed to achieve the best theoretical performance. We apply plain low-rank approximations for subdiagonal blocks of the triangular factor of the LDLT factorization, and we use the hierarchical off-diagonal low-rank (HODLR) format for approximation of the diagonal blocks. A similar approach was used by Glinskiy et al. (2017), but MPI optimization in that paper was targeted to minimize the peak memory consumption. In the current paper, the goal is performance improvement.

This paper is organized as follows: The first section discusses the setting of the boundary value problem for the Helmholtz equation, use of perfectly matched layers (PMLs) to diminish wave reflections from the boundaries, and an FD approximation of the boundary value problem. We describe a procedure for minimization of numerical dispersion and use it to construct the optimal 27-point stencil. Compared with the classic seven-point stencil, the optimization results in a $1.5 \times$ decrease of the number of points per wavelength needed to keep the numerical dispersion error within $1 \%$. In 3 D , this decrease leads to a $3 \times$ reduction of the coefficient matrix size and an even bigger factor for the performance. We describe a special ordering of the grid points defined by the ND method, which results in a special structure of the matrix.

Factorization and solving steps are the main ingredients of our direct method of solving systems of linear equations. The factorization process includes compression of matrix blocks based on the low-rank approximation. The second section is devoted to details of the software implementation. In the last section, results of the numerical experiments are given. Using two realistic subsurface models, we demonstrate the accuracy of the solver, its MPI scalability, and performance. We compare the solver with the industrial time-domain solver provided by the Seiscope consortium, and we demonstrate that both solvers complement each other.

## METHODS

## An FD scheme for the Helmholtz equation

In the frequency domain, acoustic wave propagation in $R^{3}$ is governed by the Helmholtz equation:

$$
\begin{equation*}
\Delta u(x)+k^{2}(x) u(x)=f(x) \tag{1}
\end{equation*}
$$

where $\Delta$ is the Laplace operator, $u(x)=u\left(x_{1}, x_{2}, x_{3}\right)$ is the unknown function (pressure), $k(x)=\omega / c(x)$ is the wavenumber, $\omega$ is the angular temporal frequency, and $c(x)$ is the sound velocity at point $x$ (To avoid cumbersomeness in our formulas, we use $x$ along with triples $\left(x_{1}, x_{2}, x_{3}\right)$ and $(x, y, z)$ as notations for the independent spatial variables. It should be clear from the context which notation is used, and we hope that such a shortcut would not confuse the reader.). The pressure is excited by a source represented by function $f(x)$ in the right side. For a seismic shot located at source point $x^{s}=\left(x_{1}^{s}, x_{2}^{s}, x_{3}^{s}\right)$, the right side has a form of the delta function:

$$
\begin{equation*}
f(x)=\delta\left(x-x^{s}\right) \tag{2}
\end{equation*}
$$

If equation 1 is considered in an unbounded domain, for unique resolvability some condition (e.g., radiation condition; Vainberg, 1966) should be posed at infinity.

In practice, the solution is to be found in some parallelepiped $D$ where the velocity function is given and the source point is located. To reduce the influence on the solution of the limited computational domain, we apply the PML technique (see, e.g., Berenger, 1994; Bermúdez et al., 2007). For this purpose, $D$ is immersed in some bigger parallelepiped $\tilde{D}$ as shown schematically in Figure 1. If freesurface boundary conditions need to be imposed, the respective faces of $D$ and $\tilde{D}$ lie in the same plane. The differential equation is extended to $\tilde{D}$ in the form of the second-order partial differential equation:

$$
\begin{equation*}
\sum_{i=1}^{3} \alpha_{i}\left(x_{i}\right) \frac{\partial}{\partial x_{i}}\left(\alpha_{i}\left(x_{i}\right) \frac{\partial u}{\partial x_{i}}\right)+k^{2}(x) u(\mathrm{x})=f(x), x \in \tilde{D} \tag{3}
\end{equation*}
$$

where $k(x)$ is smoothly extended to $D_{\text {PML }}$ and $f(x)$ is extended by zero. Functions

$$
\begin{equation*}
\alpha_{i}\left(x_{i}\right)=\frac{i \omega}{i \omega+d_{i}\left(x_{i}\right)} \tag{4}
\end{equation*}
$$

are complex-valued defined via damping functions $d_{i}\left(x_{i}\right)$, which are zeros for points within $D$. Respectively, $\alpha_{i}\left(x_{i}\right)$ is equal to one in $D$, and equation 3 coincides with equation 1 . To provide attenuation of the solution, damping functions are positive within the respective PMLs. To maintain the symmetry of the coefficient matrix in the FD approximation, we divide the partial differential equation 3 by $\alpha(x)=\alpha_{1}\left(x_{1}\right) \alpha_{2}\left(x_{2}\right) \alpha_{3}\left(x_{3}\right)$ and obtain

$$
\begin{equation*}
\sum_{i=1}^{3} \frac{\alpha_{i}\left(x_{i}\right)}{\alpha(x)} \frac{\partial}{\partial x_{i}}\left(\alpha_{i}\left(x_{i}\right) \frac{\partial u}{\partial x_{i}}\right)+\frac{k^{2}(x)}{\alpha(x)} u(x)=\frac{f(x)}{\alpha(x)} . \tag{5}
\end{equation*}
$$

On the boundary of domain $\tilde{D}$, the homogeneous Dirichlet condition is posed

$$
\begin{equation*}
\left.u(x)\right|_{x \in \partial \tilde{D}}=0 \tag{6}
\end{equation*}
$$

Differential equation 5 or 3 along with boundary condition 6 is the boundary value problem used for numerical simulation of acoustic waves in the frequency domain.
For numerical solution of this boundary value problem, we first introduce a rectangular equidistant grid in $\tilde{D}$. The grid points $\left(x_{i}, y_{j}, z_{k}\right)$ are defined as multiples of integer triplets $(i, j, k)$ with grid step $h$. A discrete analog $u_{i j k}$ of the unknown function $u(x, y, z)$ is defined as $u_{i j k}=u\left(x_{i}, y_{j}, z_{k}\right)$, and discretization $f_{i j k}$ of the right side $f(x) / \alpha(x)$ is similarly done.
Differential equation 5 is approximated using FD schemes. To reduce the impact of numerical dispersion, one can use FD schemes of a higher order of approximation (Dablain, 1986; Liu and Sen, 2011). Such schemes use wider stencils, which may result in cost increases at the matrix factorization stage.

To approximate the differential equation at point $\left(x_{i}, y_{j}, z_{k}\right)$, we use a 27 -point FD approximation of second-order accuracy

$$
\begin{equation*}
\sum_{i_{0}=-1}^{1} \sum_{j_{0}=-1}^{1} \sum_{k_{0}=-1}^{1} \beta_{i, i_{0}, j, j_{0}, k, k_{0}} u_{i+i_{0}, j+j_{0}, k+k_{0}}=f_{i j k} \tag{7}
\end{equation*}
$$

to boundary value problems 5 and 6 . To define coefficients $\beta_{i, i_{0}, j, j_{0}, k, k_{0}}$, we represent the left side of equation 7 as a linear combination of seven terms with the weight coefficients

$$
\begin{align*}
& \gamma_{1} \Delta_{1} u_{i j k}+\gamma_{2} \Delta_{2} u_{i j k}+\gamma_{3} \Delta_{3} u_{i j k} \\
& +k^{2}\left(x_{i}, y_{j}, z_{k}\right)\left(w_{1} u_{i j k}+\frac{w_{2}}{6} \sum_{(2)} u_{i_{1}, j_{1}, k_{1}}\right. \\
& \left.\quad+\frac{w_{3}}{12} \sum_{(3)} u_{i_{1}, j_{1}, k_{1}}+\frac{w_{2}}{8} \sum_{(4)} u_{i_{1}, j_{1}, k_{1}}\right) \tag{8}
\end{align*}
$$

Operator $\Delta_{j},(j=1,2,3)$ in sum 8 is a second-order-accuracy FD approximation of the Laplace operator represented by the stencil marked $(j)$ in Figure 2. In fact, in Figure 2 only $z$-parts of the respective stencils are depicted. Obviously, to get an approximation of the Laplace operator, the sum of weights $\gamma_{j}$ should be equal to one. Likewise, sums in parentheses in expression 8 are the second-orderaccuracy approximations of $u\left(x_{i}, y_{j}, z_{k}\right)$ that correspond to the stencils depicted in Figure 3. In these stencils, the red ball denotes the central point and does not participate in the approximations. The respective weights should also sum up to one.
For FD operator 8, the dispersion analysis gives an explicit expression $c_{p h}\left(\varphi, \theta, G ; \gamma_{1}, \gamma_{2}, \gamma_{3}, w_{1}, w_{2}, w_{3}, w_{4}\right)$ for the phase velocity of a plane wave propagating in a direction defined by spherical


Figure 1. Schematic view of domain $\tilde{D}=D \cup D_{\text {PML }}$. In $D_{\text {PML }}$, the solution is attenuated by introducing complex valued functions $\alpha_{i}\left(x_{i}\right)$ in equation 3. Homogeneous Dirichlet conditions are posed on boundary $\partial \tilde{D}$.

(2)

(3)


Figure 2. Three FD stencils for the approximation of the partial derivative $u_{z z}$.
angles $\varphi, \theta$. Here, $G=\omega h / 2 \pi c$ is the reciprocal of the number of grid points per wavelength. The function

$$
\begin{align*}
& \Psi\left(\gamma_{1}, \gamma_{2}, \gamma_{3}, w_{1}, w_{2} w_{3} w_{4}\right) \\
& \quad=\iiint\left|\frac{c_{p h}\left(\varphi, \vartheta, G ; \gamma_{1}, \gamma_{2}, \gamma_{3}, w_{1}, w_{2} w_{3} w_{4}\right)}{c}-1\right|^{2} d \varphi d \theta d G \tag{9}
\end{align*}
$$

represents the squared dispersion error averaged with respect to the directions and values of $G$. The domain of integration in integral 9 is defined as
$D=\left\{(\varphi, \theta, G): 0 \leq \varphi \leq 2 \pi,-\frac{\pi}{2} \leq \theta \leq \frac{\pi}{2}, G_{\min } \leq G \leq G_{\max }\right\}$.

Minimizing function 9 with respect to the weights, we find the optimal values. In Table 1, one can find the optimal weights for the range of frequencies $4-10 \mathrm{~Hz}$.


Figure 3. Four stencils for the approximation of $k^{2}(x) u(x)$.

Table 1. Values of the optimal parameters for stencil 8.

| $\gamma_{1}$ | 0.6558 | $w_{1}$ | 0.5756 |
| :--- | ---: | ---: | ---: |
| $\gamma_{2}$ | 0.2945 | $w_{2}$ | 0.1874 |
| $\gamma_{3}$ | 0.0497 | $w_{3}$ | 0.3588 |
|  |  | $w_{4}$ | -0.1219 |



Figure 4. For 50 randomly chosen directions of propagation of plane waves, dispersion error curves are plotted as functions of the reciprocal of the number of grid points per wavelength. The FD approximation is defined by operator 8 with optimal parameters from Table 1.

Dispersion curves for 50 plane waves propagating in randomly chosen directions are depicted in Figure 4. It is clear that eight points per wavelength is sufficient to keep the dispersion error at less than $1 \%$. To justify our choice of using a 27-point optimal stencil, we note that for the classic seven-point approximation to reach the level of $1 \%$ dispersion error, 12 points per wavelength are required. Optimization of FD stencils for minimization of the numerical dispersion is a popular topic of many papers (Jo et al., 1996; Hustedt et al., 2004; Operto et al., 2007).

## A system of linear algebraic equations

The FD approximation 7 of boundary value problems 5 and 6 takes the form of a system of linear equations:

$$
\begin{equation*}
\mathbf{A u}=\mathbf{f} \tag{11}
\end{equation*}
$$

where $\mathbf{u}$ stands for a column of values $u_{i j k}$ at grid points of the unknown function and $\mathbf{f}$ has components $f_{i j k}$. For any grid point $(i, j, k)$, there are at most 27 nonzero coefficients in equation 7. In other words, matrix $\mathbf{A}$ is sparse, complex valued (due to PMLs), and symmetric (but not Hermitian). Coefficients $\beta_{i, 0, j, 0, k, 0}$ come to the diagonal of the matrix, whereas the locations of remaining 26 nonzero coefficients depend on the ordering of the grid points.

To order the grid points, we use the ND algorithm characterized by a substantial reduction in memory consumption (George, 1973; George et al., 1994; Davis et al., 2016) compared with conventional ordering. For convenience, we illustrate the ND algorithm in Figure 5 in which the computational domain is depicted as a parallelepiped. The grid points are not shown for simplicity. The algorithm deals with subsets of the finite set of grid points, but, for visibility, we operate with geometric objects and terms. We sequentially divide the domain into subdomains using separators until the sizes of the subdomains reach a computationally manageable level.

At the first step, in Figure 5a, the separator is shown as a horizontal gray plane labeled 15 . The plane passes through the grid points and divides the parallelepiped into approximately equal subdomains. The entire set of grid points disjoins into three subsets: the upper subdomain (the first subset), the lower subdomain (the second subset), and the separator itself (the third subset).

Let us sequentially order the grid points (subset one, subset two, and subset three). At this point, the ordering within the subsets is not specified.

A stencil centered at some grid point $(i, j, k)$ connects the center with 26 neighbors, i.e., points whose grid coordinates differ from the center coordinates by not more than one. Obviously, for a center from the upper subdomain, the neighbors belong to the upper subdomain, or, possibly to the separator. From this observation, we get the following block structure

$$
\mathbf{A}=\left(\begin{array}{ccc}
\mathbf{A}_{\mathrm{uu}} & 0 & \mathbf{A}_{\mathrm{u}, \mathrm{~s}_{15}}  \tag{12}\\
0 & \mathbf{A}_{\mathrm{ll}} & \mathbf{A}_{1, \mathrm{~s}_{15}} \\
\mathbf{A}_{\mathrm{s}_{15}, \mathrm{u}} & \mathbf{A}_{\mathrm{s}_{15}, \mathrm{l}} & \mathbf{A}_{\mathrm{s}_{15}, \mathrm{~s}_{15}}
\end{array}\right),
$$

of the coefficient matrix.

The process can be repeated recursively. In Figure 5b, a vertical gray plane partitions the subdomains. This vertical plane contains two separators 13 and 14 . We order the grid points so that points from separators 13,14 , and 15 are put at the end of the list and grid points from four subdomains are arranged first, and we have a block $7 \times 7$ matrix

In this formula, matrix block indices $1,2,3$, and 4 indicate the grid points from the subdomains, whereas $\mathrm{s}_{13}, \mathrm{~s}_{14}$, and $\mathrm{s}_{15}$ relate to the points from the respective separators.

The third-level separators 9, 10, 11, and 12 are shown in Figure 5c. Every higher level separator partitions the subdomains obtained at the previous level into two smaller subdomains and the separator itself. So, $k$ steps of such process produce $2^{k}$ subdomains and $2^{k}-1$ separators. The grid points are ordered so that separators in reverse ordering follow the points from the subdomains. The coefficient matrix becomes a block $\left(2^{k+1}-1\right) \times\left(2^{k+1}-1\right)$ matrix. For $k=3$, the structure is depicted in Figure 6a.

## LDLT factorization with compression

To solve the system of linear equations 8 with a sparse complex symmetric coefficient matrix by a direct method, LDLT factorization is first applied:

$$
\begin{equation*}
\hat{\mathbf{A}}=\mathbf{P} \cdot \mathbf{A} \cdot \mathbf{P}^{\mathbf{t}}=\mathbf{L} \cdot \mathbf{D} \cdot \mathbf{L}^{\mathbf{t}} \tag{14}
\end{equation*}
$$

where $\mathbf{P}$ denotes a permutation matrix that appears due to pivoting, $\mathbf{D}$ is a block-diagonal matrix with blocks of size one or two, $\mathbf{L}$ is a lower triangular matrix with units on the diagonal, and " $t$ " indicates the transposed of the matrix. Factorization 14 is considered as a variant of Gauss elimination (see, e.g., Bunch et al., 1976; Duff et al., 1979; Duff and Reid, 1983; Schenk and Gärtner, 2006).

Let the matrix be represented in a block form as

$$
\hat{\mathbf{A}}=\left(\begin{array}{cccc}
\mathbf{A}_{11} & \mathbf{A}_{12} & \cdots & \mathbf{A}_{1 K}  \tag{15}\\
\mathbf{A}_{21} & \mathbf{A}_{22} & \cdots & \mathbf{A}_{2 K} \\
\vdots & \vdots & \ddots & \vdots \\
\mathbf{A}_{K 1} & \mathbf{A}_{K 2} & \cdots & \mathbf{A}_{K K}
\end{array}\right)
$$

with $n_{i} \times n_{j}$-blocks $\mathbf{A}_{i j}$. Diagonal blocks $\mathbf{A}_{i i}$ are square though may have different orders $n_{i}$. The block structure of matrices $\mathbf{L}$ and $\mathbf{D}$ is as follows:


Figure 5. A schematic view of the ND algorithm comprising (a) the first-level separator 15 partitioning the domain into the upper and the lower subdomains, (b) second-level separators 13 and 14 partitioning the domain into four subdomains, and (c) further partitioning by third-level separators $9,10,11$, and 12 resulting in eight subdomains.


Figure 6. The $\mathbf{L}$ factor structure showing (a) the original and (b) the low-rank approximation. The fill-in is shown by the gray shading. See the respective ET in Figure 7.
phenomenon). The phenomenon can be easily understood from equation 19 for Schur updates. The gray shading in Figure 6a is used to illustrate the "fill-in factor" (ratio of the number of nonzero elements in the block to the total number of the block elements). The fully zero blocks are white colored, and those in black mean the fill-in factor is close to one. As illustrated in Figure 6a, blocks on the lower and right side of matrix $\mathbf{L}$ have bigger fill-in factors and may be larger in size.

Figure 6 depicts the structure of factor $\mathbf{L}$, but to some extent, this also describes the lower triangle of matrix $\hat{\mathbf{A}}$ (the upper triangle can be obtained by reflecting the picture with respect to the main diagonal). The elimination tree (ET) for matrix $\hat{\mathbf{A}}$ has a binary structure (Figure 7). Recall that for the LDLT factorization, dependencies among panels are indicated by the ET edges, and these dependencies denote the directions of how updates move from a "son" to its "father," "grandfather" and beyond up to the tree "root" along the tree branches. For example, Figure 7 shows that panel 13 depends on panels 9 and 10 , which in turn depend on $1,2,3$, and 4 . Panel 13 receives updates not only from its "sons" 9 and 10 but also from "grandsons" $1,2,3$, and 4 . These updates are independent of 9 and 10 - this additional parallelism can be used for performance optimization.

Fortunately, for the FD approximation of the Helmholtz equation, matrix $L$ blocks lying below the diagonal possess data sparsity property; i.e., they are low rank. $M \times N$-matrix $\mathbf{B}$ is called data sparse if it can be approximated by a matrix $\mathbf{B}^{r}$ of rank $r \ll \min (M, N)$. Such an approximation can be written as

$$
\begin{equation*}
\mathbf{B}=\mathbf{B}^{r}+\boldsymbol{\delta} \mathbf{B}=\mathbf{U} \cdot \mathbf{V}^{t}+\boldsymbol{\delta} \mathbf{B},\|\delta \mathbf{B}\|<\varepsilon\|\mathbf{B}\| \tag{21}
\end{equation*}
$$


#### Abstract

Algorithm 1. Pseudocode of a block version of LDLT factorization. The sparsity of the matrix is taken into account by skipping values of indices $i$ and $p$ that correspond to fully zero blocks in equations 17,19 , and 20.


$$
\begin{aligned}
& \text { for } j=1: K \\
& \quad \text { for } p=1: j-1
\end{aligned}
$$

$$
\begin{equation*}
\mathbf{A}_{j j}=\mathbf{A}_{j j}-\mathbf{L}_{j p} \mathbf{D}_{p} \mathbf{L}_{j p}^{t} \tag{17}
\end{equation*}
$$

end

$$
\begin{equation*}
\left[\mathbf{L}_{j j}, \mathbf{D}_{j}, \mathbf{P}_{j}\right]=\operatorname{FLDLT}\left(\mathbf{A}_{j j}\right) \tag{18}
\end{equation*}
$$

for $i=j+1: K$
for $p=1: j-1$

$$
\begin{equation*}
\mathbf{A}_{i j}=\mathbf{A}_{i j}-\mathbf{L}_{i p} \mathbf{D}_{p} \mathbf{L}_{j p}^{t} \tag{19}
\end{equation*}
$$

end

$$
\begin{equation*}
\mathbf{L}_{i j}=\mathbf{A}_{i j} \mathbf{P}_{j}^{t} \mathbf{L}_{j j}^{-t} \mathbf{D}_{j}^{-1} \tag{20}
\end{equation*}
$$

end
end
with $M \times r$-matrix $\mathbf{U}, N \times r$-matrix $\mathbf{V}$, and $\varepsilon>0$ being some threshold parameter. Matrices $\mathbf{U}$ and $\mathbf{V}$ are stored instead of $\mathbf{B}$. Approximation 21 is used for data compression in the LDLT factorization algorithm we use; the lower the sum $r / M+r / N$, the bigger is the compression effect. The compression also helps to reduce the floating point operations count.

The low-rank approximation in equation 21 provides a foundation for data compression. For the spectral matrix norm or Frobenius norm, the truncated singular-value decomposition (SVD) (see Golub and Loan, 1996; Godunov et al., 2013) can be used to solve the minimization problem:

$$
\begin{equation*}
\mathbf{B}^{r}=\arg \min _{\operatorname{rank} \mathbf{X}=r}\|\mathbf{B}-\mathbf{X}\|^{2} \tag{22}
\end{equation*}
$$

The solution can be considered as the optimal approximation of matrix $\mathbf{B}$ by rank $r$ matrices. SVD provides a constructive solution $\mathbf{B}^{r}$ to the approximation problem 22 in the form of matrices $\mathbf{U}$ and $\mathbf{V}$ such that $\mathbf{B}^{r}=\mathbf{U} \cdot \mathbf{V}^{t}$ (see equation 21). The SVD algorithm is numerically stable but time consuming. Many alternative approaches can be found in the literature. We use an approach based on randomized sampling (Martinsson and Voronin, 2011) because of its high performance and robustness.

Data sparsity in matrix $\mathbf{L}$ blocks is schematically illustrated in Figure 6b in which some blocks are shown as containing two narrow black blocks schematically representing respective matrices $\mathbf{U}$ and $\mathbf{V}^{t}$ that are stored instead of the block they approximate. Usually these blocks are dense, which is why they are black in the figure. Obviously, due to some overhead, approximation 21 of matrix $\mathbf{L}$ blocks makes sense for high enough fill-in factors (Figure 6b), whereas some pale-gray blocks remain "as is."

There is one more possibility to compress matrix L: Diagonal blocks of this matrix in Figure 6b can be approximated by HODLR matrices (Aminfar et al., 2016; Glinskiy et al., 2017; Kostin et al., 2017). To get this structure, the respective diagonal blocks of $\mathbf{L}$ are represented as $2 \times 2$ block matrices with subblocks of half-size and approximation 21 is applied to the off-diagonal subblock. In our solver, this procedure is applied recursively to the diagonal blocks several times until a small size is reached.

## Solving step

Having defined factors in factorization 11, one can solve the system of linear equation 14 by inverting triangular matrices


Figure 7. The ET corresponding to the matrix structure in Figure 6. The tree nodes correspond to the panels in Figure 6. The dependence of panel node 13 on nodes 9 and 10 is illustrated in Figure 6a as in the respective substructure of the matrix.

$$
\begin{equation*}
\mathbf{w}=\mathbf{D}^{-1} \mathbf{L}^{-1} \mathbf{P} \mathbf{f} ; \quad \mathbf{u}=\mathbf{P}^{t} \mathbf{L}^{-t} \mathbf{w} \tag{23}
\end{equation*}
$$

In a case of several right sides, vectors $\mathbf{f}, \mathbf{y}, \mathbf{z}, \mathbf{u}$ should be treated as matrices comprised of the respective column vectors, and computations 23 can be parallelized with respect to the columns.

During computation of factorization 14 , matrix $\mathbf{L}$ is compressed to reduce memory consumption. The compressibility factor depends on threshold $\varepsilon$, frequency $\omega$, and some other parameters; in our experiments, the compressibility factor varied from 3 to 7. Due to compression, equation 14 is no longer valid and should be replaced with

$$
\begin{equation*}
\hat{\mathbf{A}}=\mathbf{L} \cdot \mathbf{D} \cdot \mathbf{L}^{t}+\boldsymbol{\delta} \hat{\mathbf{A}} \tag{24}
\end{equation*}
$$

where $\boldsymbol{\delta} \hat{\mathbf{A}}$ denotes a perturbation due to compression errors. The solution obtained by formula 21 may become inaccurate. To improve the accuracy, one can carry out a few steps of the iterative refinement (Wilkinson, 1964). Having computed factorization 14, the overhead due to inclusion of the iterative refinement in the solving process is low.

Provided that the norm of perturbation $\boldsymbol{\delta} \hat{\mathbf{A}}$ is small enough and the coefficient matrix is "not-too-badly-conditioned," the iterative refinement converges. The threshold $\varepsilon$ directly impacts the compressibility factor - the bigger its value, the bigger the compressibility factor, suggesting that the threshold value should be increased. However, a larger threshold value implies an increase in the norm of perturbation $\boldsymbol{\delta} \hat{\mathbf{A}}$ in equation 24 , which may lead to a loss of convergence of the iterative refinement, or at least it may require more iterations to converge. Therefore, unjustifiably large threshold values should be avoided. Due to this trade-off, in practice, the threshold value is established by trial and error.

## SOFTWARE IMPLEMENTATION DETAILS

Some standard functionality needed for our solver is taken from the $\operatorname{Intel}(\mathrm{R})$ Math Kernel Library (Intel(R) MKL, 2019). This library provides highly optimized Basic Linear Algebra Subroutines (BLAS) (BLAS, 2017) and Linear Algebra Package (LAPACK) (LAPACK, 2017) functions. In our solver, we use Open Multi-Processing (OpenMP) parallelization implicitly by calling threaded versions of BLAS and LAPACK functions. This is the only way of using OpenMP threading in the solver; there is no explicit OpenMP parallelization in it.

Our solver is targeted for distributed memory systems (clusters) using MPI. The binary structure of the ET of the LDLT factorization algorithm provides good MPI parallelization opportunities. Figure 8 shows mapping of the ET structure on compute nodes for the case of four nodes. Each subtree (denoted by the dashed ellipses) is assigned to a particular compute node. The nodes are assumed to have big enough RAM to contain the respective input data and results. It is clear that up to the level 3 (in Figure 8, this level is marked with a dashed line passing the ET nodes) of the ET, the compute nodes process their portions of data independently and there is no data exchange between them.

However, this kind of parallelism decreases while the process traverses up along the tree, and at a particular level of the tree, the number of nodes becomes less than the number of cluster nodes. Moreover, while moving up along the ET, fill-in factors of the respective panels increase with increasing memory requirements and numbers of floating point operations.

Upon reaching the level in which the ET width becomes less than the number of cluster nodes, finer granulation of the matrix being processed is applied. In other words, the panels are subdivided into subpanels, which are assigned to separate cluster nodes. Unfortunately, such a subdivision does not provide a binary structure of the ET suitable for parallelization. However, it allows delegating computation of some Schur updates to the lower part of the tree, thus improving the parallelism.

One of the costly operations in the LDLT factorization is the Schur update

$$
\begin{equation*}
\mathbf{A}_{i j}=\mathbf{A}_{i j}-\sum_{p} \mathbf{L}_{i p} \mathbf{D}_{p} \mathbf{L}_{j p}^{t} \tag{25}
\end{equation*}
$$

(see the $p$-loop in Algorithm 1). For $j$ th node of the ET, equation 25 means gathering this node updates from sons, grandsons, etc. The number of summands in this formula may be rather large, and matrices $\mathbf{L}_{i p}$ are given factorized in products of "thin" matrices and they are distributed along different cluster nodes. To apply the Schur update 25 , we use different strategies depending on the position of the ET node with respect to the dashed line.

For the ET nodes lying above the dashed line, or, equivalently, for the values $j$, we split the total number of updates (i.e., indices $p$ ) into pairs. Each pair is summed up on a separate cluster node. The process is repeated several times until all updates are summed up. Finally, the sum of all updates is added to block $\mathbf{A}_{i j}$. The number of steps is proportional to the logarithm of the total number of updates.

For the ET nodes below the dashed line, all operations are performed on the same compute node but there are also some peculiarities. As was already mentioned above, nonzero subdiagonal blocks in the leftmost panels of the triangular factor are stored as is, without applying the low-rank approximation. This is beneficial from the performance and memory consumption points of view. So, to compute Schur update 25 with blocks in the conventional format, standard BLAS functionality (with some modifications due to the sparsity of the blocks) is needed.

Starting from some level, the low-rank approximation is applied to the blocks and the low-rank arithmetic is switched on. Let


Figure 8. Mapping of the ET to the cluster nodes. The dashed ellipses denote the ET nodes combined together and assigned to the same cluster node.
$m_{1} \times n$-block $\mathbf{L}_{i p}$ and $m_{2} \times n$-block $\mathbf{L}_{j p}$ be represented by low-rank approximants $\mathbf{U}_{i p} \cdot \mathbf{V}_{i p}^{t}$ and $\mathbf{U}_{j p} \cdot \mathbf{V}_{j p}^{t}$, where $\mathbf{U}_{i p}$ and $\mathbf{V}_{i p}$ have sizes $m_{1} \times r_{1}$ and $n \times r_{1}$, whereas $\mathbf{U}_{j p}$ and $\mathbf{V}_{j p}$ have sizes $m_{2} \times r_{2}$ and $n \times r_{2}$, respectively. Assuming $r_{1} \leq r_{2}$, update $\mathbf{L}_{i p} \mathbf{D}_{p} \mathbf{L}_{j p}^{t}$ becomes a product of $m_{1} \times r_{1}$-matrix $\mathbf{U}_{i p}$ and $r_{1} \times m_{2}$-matrix $\left(\mathbf{U}_{j p}\left(\mathbf{V}_{j p}^{t}\left(\mathbf{D}_{p}^{t} \mathbf{V}_{i p}\right)\right)^{t}\right)^{t}$ (if the opposite inequality $r_{1}>r_{2}$ takes place, matrix $\mathbf{L}_{i p} \mathbf{D}_{p} \mathbf{L}_{j p}^{t}$ is decomposed in a product of $m_{1} \times r_{2}-$ matrix and $r_{2} \times m_{2}$-matrix).

Summing up rank $r_{1}$ update $\mathbf{U}_{1} \cdot \mathbf{V}_{1}^{t}$ and rank $r_{2}$ update $\mathbf{U}_{2} \cdot \boldsymbol{V}_{2}^{t}$ can be obtained by the concatenation of matrices

$$
\mathbf{U}_{1} \cdot \mathbf{V}_{1}^{t}+\mathbf{U}_{2} \cdot \mathbf{V}_{2}^{t}=\left(\begin{array}{lll}
\mathbf{U}_{1} & \mathbf{U}_{2}
\end{array}\right)\left(\begin{array}{ll}
\mathbf{V}_{1} & \mathbf{V}_{2} \tag{26}
\end{array}\right)^{t}
$$

The result has the desired form but probably the representation is not optimal - matrices $\left(\begin{array}{ll}\mathbf{U}_{1} & \mathbf{U}_{2}\end{array}\right)$ and $\left(\begin{array}{ll}\mathbf{V}_{1} & \mathbf{V}_{2}\end{array}\right)$ have $r_{1}+r_{2}$ columns that may be greater than their ranks. In other words, the product needs to be recompressed

$$
\left(\begin{array}{lll}
\mathbf{U}_{1} & \mathbf{U}_{2}
\end{array}\right)\left(\begin{array}{ll}
\mathbf{V}_{1} & \left.\mathbf{V}_{2}\right)^{t}=\mathbf{U} \mathbf{V}^{t} . \tag{27}
\end{array}\right.
$$

This is the same problem as the compression problem.

## 3D NUMERICAL EXPERIMENTS

In this section, we describe a few numerical experiments to validate our implementation. The experiments were performed on the Linux supercomputer Shaheen II (n.d.) at the King Abdullah University of Science and Technology (KAUST, Saudi Arabia). Top 500 List (2018) Linpack performance of this cluster is 5.5 Pflops/s. Each compute node is supplied with two Intel Xeon CPU processors E5-2698 v3 @ 2.3 GHz ( 32 cores in total) and 128 GB RAM; this configuration provides theoretical peak performance of 1.2 Tflops/s per node.

The threshold parameter was taken from the range $10^{-5}$ to $10^{-4}$. We used one MPI process per cluster node. The stopping criterion for the iterative refinement was set to achieve the following inequality:

$$
\begin{equation*}
\frac{\|\mathbf{A u}-\mathbf{f}\|}{\|\mathbf{f}\|}<10^{-5} . \tag{28}
\end{equation*}
$$

## Models

To demonstrate the accuracy of the method, first we compare the numerical solution for a homogeneous model to the analytical solution. For a synthetic model with the velocity being linearly dependent on the depth, we compare solutions obtained with different grid steps.


Figure 9. Cross sections of the TZ velocity model.

The second subsurface model is a marine transition zone model referred below as the TZ model with dimensions of $18,000 \times 23,500 \times 7000 \mathrm{~m}$. In this model, the velocity varies between 1042 and $7626 \mathrm{~m} / \mathrm{s}$ (Figure 9). FD discretization with a grid size of 30 m results in a system of $1.3 \times 10^{8}$ linear equations.

As a third subsurface model, we use a portion of the SEG overthrust (referred to below as OT) model (Aminzadeh et al., 1997) with dimensions of $9870 \times 9870 \times 4620 \mathrm{~m}$. In this model, the velocity varies between 2286 and $6000 \mathrm{~m} / \mathrm{s}$. FD discretization with a grid of 30 m in all directions results in a system of $2.4 \times 10^{7}$ linear equations.

## Accuracy

First, we benchmark numerical solution of the Helmholtz equation:

$$
\begin{equation*}
\left(\Delta+k^{2}\right) v=-\delta(x) \tag{29}
\end{equation*}
$$

with the analytical solution available for a homogeneous medium

$$
\begin{equation*}
v(x, y, z)=\frac{e^{-i k \sqrt{x^{2}+y^{2}+z^{2}}}}{4 \pi \sqrt{x^{2}+y^{2}+z^{2}}} \tag{30}
\end{equation*}
$$

We compute numerical solutions for different values of the grid step. Accuracy of the approximate solution $u^{h}(x, y, z)$ is measured by metrics

$$
\begin{equation*}
\beta_{k}(r)=\frac{\left\|u^{h}(x, y, z)-v(x, y, z)\right\|_{B\left(r_{0}, r\right), k}}{\|v(x, y, z)\|_{B\left(r_{0}, r\right), k}} \tag{31}
\end{equation*}
$$

depending on the size $r$ of the domain. In definition 31, index $k$ stands for the type of the norm $(k=1,2)$ and

$$
\begin{equation*}
B\left(r_{0}, r\right)=\left\{(x, y, z) \mid r_{0} \leq \sqrt{x^{2}+y^{2}+z^{2}} \leq r\right\} \tag{32}
\end{equation*}
$$

denotes the concentric spherical layer in which the norms are computed. Parameter $r_{0}$ is used to exclude the small singularity region around the origin. Metrics $\beta_{k}(r)$ are simple relative accuracies of the solution for the respective norms. Note that the accuracy also depends on the grid step $h$.

We provide numerical results for the homogeneous model with dimensions of $5120 \times 5120 \times 2560 \mathrm{~m}$ and sound velocity $c_{0}=1280 \mathrm{~m} / \mathrm{s}$. The wavefield was excited by a monochromatic 4 Hz point source placed at the center of the model $(2560,2560$, and 1280 m ). We used three different grid steps $h=64,32$, and 16 m that correspond to 5,10 , and 20 points per wavelength, respectively. The width of the PMLs was taken to be 15 grid steps. In Figure 10, one can find graphs of functions $\beta_{k}(r)$ for the described problem setting. It is clear that halving the grid step results in a four times reduction of the relative errors $\beta_{1}$ and $\beta_{2}$, which confirms the secondorder approximation of the method.

For the second test, we used a 3D model with the velocity being a linear function of depth and investigated convergence of the solution with a decreasing grid step. For the same velocity model, we computed solutions $u^{h}, u^{h / 2}$, and $u^{h / 4}$ for steps $h, h / 2$, and $h / 4$,
respectively. The solution $u^{h / 4}$ was used as a substitute for the exact solution. In Figure 11, one can find the respective functions $\beta_{k}(r)$, which in this case provide information on the deviation of solutions $u^{h}, u^{h / 2}$ from $u^{h / 4}$. The second-order approximation is clearly confirmed for $\beta_{1}$ and $\beta_{2}$.

## Strong scaling

For a problem that can be solved with an MPI solver and parallelization on different numbers of cluster nodes, the "scaling factor" is defined as ratio $t_{\text {seq }} / t_{P}$ where runtimes $t_{\text {seq }}$ and $t_{P}$ stand for one and $P$ cluster nodes, respectively.

For the OT velocity model and different values of frequencies, we measured the scaling factors of our solver (Figure 12). The fre-quency-domain FD (FDFD) solver scales well for moderate cluster sizes, but scaling becomes worse as the number of cluster nodes increases beyond approximately four to eight nodes. In addition, scaling factors deteriorate with increasing frequency. The problem size also impacts the scaling, with better scaling for larger models as illustrated in Figure 13 in which scaling factors for the OT model and the TZ models are compared.
Scaling behavior can be explained as follows. The ET is less balanced on lower levels than on upper levels. The amount of


Figure 10. Relative errors of numerical solutions computed for grid steps $h, h / 2$, and $h / 4(64,32$, and 16 m$)$ in a homogeneous medium as a function of the domain radius $r$.


Figure 11. Relative deviation of solutions $u^{h}, u^{h / 2}$ from $u^{h / 4}$ for a model with the velocity being a linear function of the depth.
computations defined by the dashed ellipses in Figure 8 may vary from one ellipse to another. In other words, different cluster nodes are assigned different amounts of work, becoming unbalanced. Spatial variations of the sound velocity may cause unbalancing because the numerical ranks depend on the local velocity and the frequency. At the highest level of the ET, the subtrees represent big parts of the domain that appear as results of the ND algorithm. Most likely, the "average velocity" for these big parts experiences minor variation from one part to another. However, if we move down along the tree, the granulation becomes finer and variation of the velocity becomes higher. This leads to the larger difference of numerical ranks and the associated unbalancing of the cluster nodes.
A similar cause of unbalancing also appears at lower levels of the ET when the parts of the computational domain become small parallelepipeds. These parallelepipeds have faces of two different types - the faces that are parts of the outer boundary of the domain and the faces that are parallel to separators. A small parallelepiped


Figure 12. The scaling factors for the direct solver as functions of the number $P$ of cluster nodes shown for the OT velocity model.


Figure 13. Scaling factors $t_{32} / t_{P}$ for the $\mathrm{TZ}\left(1.8 \times 10^{8}\right.$ equations) and OT models ( $2.4 \times 10^{7}$ equations). Observe the higher scalability for the larger TZ model.
may have from zero to three faces that are parts of the outer boundary and from six to three faces parallel to the separators. These differences cause different amount of computations associated with the different nodes of the ET.

Solution of a problem with the TZ model requires at least 32 cluster nodes, and we cannot use the definition of the scaling factor described above. In Figure 13, we show graphs of the ratios $\frac{t_{32}}{t_{p}}$, where $t_{32}$ stands for runtime on 32 cluster nodes. Figure 13 also shows data for the smaller OT model revealing that scaling and hence performance are better for larger systems of equations using the FDFD approach.

## Comparison to a time-domain solver

We compare our solver with a time-domain FD (TDFD) solver developed by the SEISCOPE Consortium (2019). To simulate wave propagation in the time domain, the initial-boundary value problem for the wave equation is solved by TDFD. To eliminate reflection of waves from the boundaries, the domain is surrounded with PMLs. The differential equation is approximated using a fourth-order explicit FD scheme. The MPI implementation of the TDFD solver uses parallelization with respect to shots, so that one shot is assigned per single cluster compute node. No data exchange between nodes is performed because the solutions for different shots are fully independent of each other. Dependence of the TDFD solver runtime on the number of seismic shots and the number of compute nodes can be described by the formula

$$
\begin{equation*}
t_{\mathrm{TDFD}}\left(N_{\text {nodes }}, N_{\text {shots }}\right)=t_{\mathrm{TDFD}}(1,1)\left[\frac{N_{\text {shots }}}{N_{\text {nodes }}}\right] \tag{33}
\end{equation*}
$$

where $[a]$ stands for the least integer greater than or equal to $a$ and $t_{\text {TDFD }}(1,1)$ depends on some other parameters such as the problem size, the sound velocity variability, and the hardware parameters.

To get the solution in the frequency domain, the Fourier transform is applied to the time domain solution. In TDFD, the frequency spectrum of the shots contains all time frequencies, so multiple frequency components are obtained all at once.

To underline the different approaches, our method is referred to as the FDFD solver. Results of accuracy tests are provided in


Figure 14. For the OT model, the relative deviations of the solutions obtained with the TDFD and FDFD methods.

Figure 14 (the OT model) and Figure 15 (the TZ model) in terms of functions $\beta_{k}(r)$ introduced in the previous section. In Figures 14 and 15 , these functions are used to demonstrate the difference between the solution obtained with TDFD solver and the solution obtained with the FDFD solver.

For both models, solutions were computed using the parameters listed in Table 2. Note that, for each model, the solutions were computed for three different frequencies. The corresponding numbers of grid steps for the FDFD solver are shown in columns $N_{x}, N_{y}$, and $N_{z}$, whereas the total number of grid points is in column $N$, and the number of points per wavelength is listed under the "ppw" column. For the TZ model, snapshots of the real part of the solution computed along the horizontal plane at a depth of one grid step are shown in Figure 16.

The time-domain solution obtained via the TDFD solver contains components that correspond to all three values of the frequency. So, the solver was run once with a grid step of 30 m . The modeling time interval was [ $0,10 \mathrm{~s}$ ].

For particular pairs of values ( $N_{\text {nodes }}, N_{\text {shots }}$ ), the measured runtimes (in s) are provided in Tables 3 and 4 in the form of fractions, where the runtime $t_{\text {TDFD }}=t_{\text {TDFD }}\left(N_{\text {nodes }}, N_{\text {shots }}\right)$ is put as the numerator, whereas $t_{\mathrm{FDFD}}\left(N_{\text {nodes }}, N_{\text {shots }}\right)$ is put as the denominator. Note that the value of time for FDFD comprises the sum of the runtime to solve the problem for all three frequencies. Ideal scalability of the TDFD solver with respect to $N_{\text {nodes }}$ was assumed to fill in the table instead of repeating multiple shot.

Tables 3 and 4 summarize computation times of the two solvers in terms of $N_{\text {shots }}$ and $N_{\text {nodes }}$ (Figure 17). For few shots, TDFD


Figure 15. The same as Figure 14 but for the TZ model.

Table 2. Parameters of the numerical experiments.

| Model | $\nu(\mathrm{Hz})$ | $h(\mathrm{~m})$ | ppw | $N_{x}$ | $N_{y}$ | $N_{z}$ | $N=N_{x} N_{y} N_{z}$ |
| :--- | :---: | :---: | :---: | :---: | :---: | ---: | ---: |
| OT | 5 | 90 | 5.1 | 110 | 110 | 52 | 629,200 |
|  | 7 | 60 | 5.4 | 165 | 165 | 78 | $2,123,550$ |
|  | 15 | 30 | 5.1 | 330 | 330 | 155 | $16,879,500$ |
| TZ | 2 | 90 | 5.8 | 200 | 261 | 77 | $4,019,400$ |
|  | 3 | 60 | 5.8 | 300 | 391 | 116 | $13,606,800$ |
|  | 7 | 30 | 5.0 | 600 | 781 | 231 | $108,246,600$ |

solvers are relatively efficient, but for larger numbers of shots and nodes, FDFD becomes relatively more efficient per shot. The thick line (Figure 17) is the "line of equal performance" of the two solvers. For a given number of $N_{\text {nodes }}$, this line defines the number of shots that is "big enough" to fully reap the benefits of the FDFD solver and reach or exceed the numerical performance of TDFD solvers. For example, for a problem with a comparatively small number of shots that has to be solved on a particular cluster (a point close to point C of line segment CD ), the fastest way would be using the TDFD solver. As the number of shots increases (the point moves to D along CD ), starting from some number of shots, the FDFD direct solver becomes faster than the TDFD solver. Another way to look at it is to fix the number of shots, such as defined by horizontal line AB . Then moving from point A toward point B along horizontal segment $A B$ means solving the same problem on clusters with increasing number of nodes. For a fixed number of shots, the FDFD direct solver would usually be faster for comparatively smaller clusters, but on bigger clusters the TDFD solver performs better. Such behavior can be explained with the help of the following reasoning. For the TDFD solver, the runtime is roughly proportional to $N_{\text {shots }}$ with coefficient $t_{\text {TDFD }}(1,1) / N_{\text {nodes }}$ (see equation 33). The runtime for the FDFD solver can be described by a sum $t_{\text {fact }}\left(N_{\text {nodes }}\right)+t_{\text {solve }}\left(N_{\text {nodes }}, N_{\text {shots }}\right) \quad$ representing time for the factorization and time for the solving step correspondingly. In this sum, the first summand does not depend on the number of shots. To some extent, the dependence of the summand can be thought of as approximately linear $t_{\text {solve }}\left(N_{\text {nodes }}, N_{\text {shots }}\right) \cong t_{\text {solve }}\left(N_{\text {nodes }}, 1\right) \cdot N_{\text {shots }}$. The values of coefficients $t_{\text {TDFD }}(1,1) / N_{\text {nodes }}$ and $t_{\text {solve }}\left(N_{\text {nodes }}, 1\right)$ can be taken from Table 3 or 4 . For example, for the OT model $t_{\text {TDFD }}(1,1)=161$, $t_{\text {solve }}(16,1)=0.4$. This simple insight explains that if the number of shots increases, then FDFD starts to win from a certain value of shots.
In general, the necessary and sufficient condition for the existence of points ( $N_{\text {nodes }}, N_{\text {shots }}$ ) where the FDFD wins is the following:

$$
\begin{equation*}
\frac{t_{\mathrm{TDFD}}(1,1)}{N_{\text {nodes }}}>t_{\text {solve }}\left(N_{\text {nodes }}, 1\right) . \tag{34}
\end{equation*}
$$

To satisfy this condition, the right side of this inequality should be inversely proportional to $N_{\text {nodes }}$, or, in other words, the solving step is MPI parallelized.

## Performance: Comparison to Intel® MKL parallel direct sparse solver for clusters

To estimate the efficiency of a solver, its performance data in giga flops per second (GFlops) are used. Among other parameters (the hardware, software, number of MPI processes, etc.), these data depend on the problem (more precisely, the size of the model and some parameters that characterize the problem). To get performance data for the solver, one has to count the number of floating point operations used to compute the solution of the problem and measure the runtime. This approach works well for comparatively simple algorithms in which the flops counts can be calculated. The sparsity
of the matrix, use of the low-rank approximation, incorporating the iterative refinement in computations makes impossible rigorous counting of the floating point operations.
Comparison to another solver is an alternative approach to get an impression on the solver efficiency. We compared our solver with parallel direct sparse solver for clusters from $\operatorname{Intel}(\mathrm{R})$ MKL (Intel(R) Parallel Studio XE 2017 Update 4) referred to below as PARDISO. The comparison was made on a small homogeneous model of size $201 \times 201 \times 101$ referred to below as SH and the medium size OT model described above (see Table 5). RAM data in Table 5 are the peak memory usage per node. It is clearly seen


Figure 16. Horizontal snapshots of the real part of the solution for the TZ model. (a) Frequency $v=2 \mathrm{~Hz}$, depth 90 m and (b) $v=7 \mathrm{~Hz}$, depth 30 m .


Figure 17. Relative performance of the TDFD and FDFD solvers shown with the number of shots (vertical axis) versus the number of nodes (the horizontal axis). The thick line defines the line of equal performance.

Table 5. Performance and memory comparison of our solver and the MKL Cluster PARDISO.

| Model | Measured data | Our solver | PARDISO |
| :--- | :---: | :---: | ---: |
| SH | Factorization time (s) | 967 | 543 |
|  | RAM (GB) | 8.1 | 23 |
| OT | Factorization time (s) | 836 | 2615 |
|  | RAM (GB) | 17.9 | 49 |

Note: Data obtained on 129 compute nodes using 129 MPI processes and 20 OMP threads. Each node has 128 GB RAM. The threshold for our solver was $10^{-4}$.
that our solver has an essential advantage in memory usage but has worse performance on small problems. Starting from moderate-size problems, the solver outperforms the MKL Cluster PARDISO.

## CONCLUSION

We present a direct method for parallel solution of the acoustic wave equation in 3D heterogeneous media. To reduce memory consumption, the method uses intermediate data sparsity for compression. The compression technique is based on low-rank approximation of fill-in blocks. It is applied directly to blocks lying below the diagonal of the triangular factor. The diagonal blocks are also compressed using the HODLR format. These methods help to solve a system of approximately $10^{9}$ equations corresponding to the large velocity models of interest to geophysical exploration. The strong scalability of the solver is good for a moderate number of cluster nodes and large enough problems. This direct solver is successfully used for geophysical modeling applications, and some challenges remain to be resolved. Poor MPI scalability of the LDLT factorization for the number of compute nodes beyond 16 is one of the challenges.

Comparing the FDFD solver with a TDFD solver reveals that each method has its strengths and weaknesses. Detailed numerical experiments with real-world scenarios using a transition zone model and Shaheen II supercomputer demonstrate the existence of the line of equal performance defined in terms of number of shots and available nodes. The time-domain solver performs better for
comparatively large ratios of $N_{\text {nodes }} / N_{\text {shots }}$, whereas the direct solver is relatively more efficient for smaller ratios of $N_{\text {nodes }} / N_{\text {shots }}$.

Although generally larger number of nodes become available with time (as computing power becomes cheaper), seismic acquisition experiences significant growth in trace density per square kilometer enabled by larger number of sources and receivers. The concurrent "data tsunami" and increase in computing power may lead to diverse scenarios that the geophysical community needs to be prepared to handle. Depending on the available computing power and number of shots for a particular imaging or FWI problem, one or the other solver may be significantly more efficient. Therefore, the optimal FWI toolbox should contain both solvers and apply the most efficient for the specific scenario based on the line of equal performance.

In the future, the techniques used for solving the Helmholtz equation are planned to be extended to more complicated media (elasticity with possible anisotropy and viscosity). The current version of the solver can solve approximately $10^{9}$ equations. Additional optimization of the factorization step should allow computations for larger models in the near future.

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## DATA AND MATERIALS AVAILABILITY

Data associated with this research are confidential and cannot be released.

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